

Exchange of Message using Fourier Sine Transforms via Affine Transformation

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Abstract:

In this paper, we established an application of Fourier sine transformation in exchanging the message in a secured channel via affine cryptosystem which is helpful in digital electronics and signal processing.

Keywords - Decryption, Encryption, Fourier sine transform, Modulo function.

I. INTRODUCTION

Definition: Fourier sine transform [7] is an integral transform that are mainly applicable for signal processing or statistics. If $f(x)$ is defined for all positive values of x .
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$$F_s[f(x)] = \int_0^{\infty} f(x) \sin(ux) dx = F_s[u]$$

Inverse Fourier sine transform is given by.

$$f(x) = \frac{2}{\pi} \int_0^{\infty} F_s(u) \sin(ux) du$$

Properties of Fourier Sine Transforms:

1. Linearity Property:

$$F_s[a f(x) \pm b g(x)] = a F_s[f(x)] \pm b F_s[g(x)]$$

2. Change of Scale Property:

$$F_s[f(ax)] = \frac{1}{|a|} F_s\left(\frac{u}{a}\right)$$

3. Modularity Property of Fourier Sine Transforms:

$$F_s[F_s(f(x))] = \frac{1}{2} [F_c(u-a) - F_c(u+a)]$$

$$F_s[f(x)\cos ax] = \frac{1}{2} [F_s(u+a) + F_s(u-a)]$$

Cryptography: Cryptography is the science of using mathematics to hide the information. It allows us to store sensitive information or to transmit it over insecure network, so that it can only be read by the intended recipient.

Plaintext: Information that can be directly read.

Ciphertext: Encrypted data of plaintext is ciphertext.

Encryption: Process of converting plaintext to ciphertext.

Decryption: Process of reverting ciphertext to plain text.

Cryptography mainly classified in to 3 types:

1. Symmetric / private key cryptography uses single key for both encryption and decryption.
2. Hash functions: it is one way function which is infeasible practically to reverse the computation. These are the basic tools of modern cryptography [5].
3. Asymmetric / public key cryptography uses different keys for encryption and decryption.

In this paper, we propose a method of exchanging the message using Fourier sine transform via affine cryptosystem.

The plain text and ciphertext are broken up into message units. A message unit can be a single letter, also called monograph, a pair of letters called digraph, a triple of letters called trigraph or a block of more than 3 letters called multigraph [3].

Affine transformation is one of the types of symmetric key cryptosystems.

Affine transformation is defined as follows:

$C = f(p) = ap + b \pmod N$ where 'P' is the plaintext and 'C' is the cipher text respectively. a, b and N are positive integers and 'f' is a mapping from P to C.

The plaintext 'P' can be recovered from the given cipher text 'C' i.e.,

$$P = a^{-1}(C - b) \pmod N$$

$$P = a^{-1}C - a^{-1}b \pmod N$$

$$= k_1 + k_2$$

where $k_1 = a^{-1}$ and $k_2 = -a^{-1}b$ and a^{-1} is the inverse of 'a'.

Theorem: Given the affine map $C \equiv ap + b \pmod{N}$ where $a \in (\mathbb{Z}/N\mathbb{Z})^\times, b \in (\mathbb{Z}/N\mathbb{Z})$. The transformation gives a unique value of 'p' for a given C iff $\gcd(a, N) = 1$ and that the total number of affine transformations is given by $N \cdot \phi(N)$, where ϕ is Euler-phi function [2].

Encryption Algorithm:

Step 1: Consider the plain text MATHEMATICS and write the numerical equivalents.

Step 2: Choose 'a' and 'b' in affine transformation $ax + b \pmod{26}$ such that $\gcd(a, 26) = 1, \gcd(b, 26) = 1$

Step 3: Now consider Fourier sine transformation and substitute the obtained numerical value of cipher text for "a" in the Fourier sine transform $\int_0^\infty e^{-ax} \sin(ux) dx$ and sender sends this ciphered message to receiver.

Decryption Algorithm:

Step 1: Receiver receives the cipher text and first decrypt by using inverse Fourier sine transformation.

Step 2: By the obtained text from step 1 receiver again decrypts the cipher text using inverse affine transformation and with suitable decryption key.

Step 3: Receiver can retrieve the plain text.

EXAMPLE:

Consider the plain text "MATHEMATICS" write the numerical equivalent of each alphabet by taking $a = 7, b = 6$ in affine transformation ("ax + b"), we get

M	A	T	H	E	M	A	T	I	C	S
12	0	19	7	4	12	0	19	8	2	18
$(7x+6) \pmod{26}$										
12	6	9	3	8	12	6	9	10	20	2

Calculations of $(7x+6) \pmod{26}$:

Alphabet	Numerical equivalent	$(7x + 6) \pmod{26}$	Value
M	12	$(7(12) + 6) \pmod{26}$	12
A	0	$(7(0) + 6) \pmod{26}$	6
T	19	$(7(19) + 6) \pmod{26}$	9
H	7	$(7(7) + 6) \pmod{26}$	3
E	4	$(7(4) + 6) \pmod{26}$	8
M	12	$(7(12) + 6) \pmod{26}$	12
A	0	$(7(0) + 6) \pmod{26}$	6
T	19	$(7(19) + 6) \pmod{26}$	9
I	8	$(7(8) + 6) \pmod{26}$	10
C	2	$(7(2) + 6) \pmod{26}$	20
S	18	$(7(18) + 6) \pmod{26}$	2

Again, by using Fourier sine transform encrypt the obtained numerical values as 'a' in the Fourier sine transform

$$\int_0^\infty e^{-ax} \sin(ux) dx, \int_0^\infty e^{-12x} \sin(ux) dx, \int_0^\infty e^{-9x} \sin(ux) dx, \int_0^\infty e^{-3x} \sin(ux) dx, \int_0^\infty e^{-8x} \sin(ux) dx, \int_0^\infty e^{-12x} \sin(ux) dx, \int_0^\infty e^{-6x} \sin(ux) dx, \int_0^\infty e^{-9x} \sin(ux) dx, \int_0^\infty e^{-10x} \sin(ux) dx, \int_0^\infty e^{-20x} \sin(ux) dx, \int_0^\infty e^{-2x} \sin(ux) dx$$

and sender sends cipher text as

$$\left[\frac{12}{2(\pi^2 + 144)}, \frac{6}{2(\pi^2 + 36)}, \frac{9}{2(\pi^2 + 81)}, \frac{3}{2(\pi^2 + 9)}, \frac{8}{2(\pi^2 + 64)}, \frac{12}{2(\pi^2 + 144)}, \frac{6}{2(\pi^2 + 36)}, \frac{9}{2(\pi^2 + 81)}, \frac{10}{2(\pi^2 + 100)}, \frac{20}{2(\pi^2 + 400)}, \frac{2}{2(\pi^2 + 4)} \right]$$

We made Fourier sine transformation as public key and affine transformation as private key to decrypt the message.

Decryption:

Recipient receives the message from the sender and first decrypt by using inverse Fourier sine transformation and get decrypts as

$$[e^{-12x}, e^{-6x}, e^{-9x}, e^{-3x}, e^{-8x}, e^{-12x}, e^{-6x}, e^{-9x}, e^{-10x}, e^{-20x}, e^{-2x}]$$

Again, by using private key as affine transformation with $E^{-1}(y) = 15(y - 6) \text{ mod } 26$

y	12	6	9	3	8	12	6	9	10	20	2
y-6	6	0	3	-3	2	6	0	3	4	14	-4
15(y-6)	90	0	45	-45	30	90	0	45	60	210	-60
Mod26	12	0	19	7	4	12	0	19	8	2	18

In this manner, we can exchange the message in a secured channel, which is more secure than the symmetric key cryptosystem.

II. CONCLUSION

We can extend this encryption scheme by using Z-transformation and by using any symmetric key cryptosystem like vigner cipher etc, as private key.

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